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ANTC Seminar 29 March, 2011

Partially Recursive Acceptance Rejection

How do I roll a 7-sided die?

Repeat

Let X be roll of 10-sided die

Until $X \in \{1, \dots, 7\}$.

~~Chance~~ $P(X=4) = 1/7$

$$= \frac{1}{10} + \frac{3}{10} \cdot \frac{1}{10} + \left(\frac{3}{10}\right)^2 \left(\frac{1}{10}\right) + \dots$$

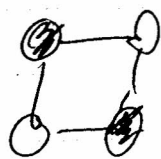
7 1st roll 8, 9, 10 1st roll 7 second roll

$$= \frac{1/10}{1 - 3/10} = \frac{1}{7}$$

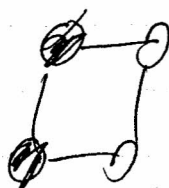
High dimensional Problem, doesn't work so well

Independent Sets

Def: An independent set of a graph is a subset of nodes where no two nodes in subset are connected by an edge.



ind set



not an ind set.

Put
up
on
Board

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Goal: • Generate samples from weighted ind sets

$$\pi_\lambda(A) = \frac{\lambda^{\#A}}{Z}$$

• Can be used to approximate Z

Applications: 1) Z is #P complete

2) Simple model of repulsive points
(cities, ~~flowers~~ trees).

Basic Acceptance rejection

- 1) Repeat
- 2) $A \leftarrow \emptyset$
- 3) For all $v \in V$
- 4) Add v to A w/ prob. $\lambda / (1 + \lambda)$
- 5) Until A is an independent set

4), 5) Treat graph like it had no edges.

⊗ ○

⊗ ○

reject

$$\begin{aligned} P(A) &= \left(\frac{\lambda}{1+\lambda} \right)^{\#A} \left(\frac{1}{1+\lambda} \right)^{\#V - \#A} \\ &= \underbrace{\left(\frac{1}{1+\lambda} \right)^{\#V}}_{\text{Constant}} \cdot \underbrace{\lambda^{\#A}}_{\propto \pi_\lambda} \end{aligned}$$

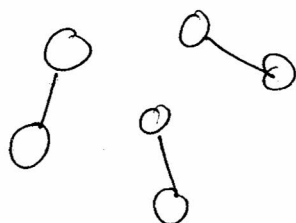
Good enough for acceptance rejection.

Problem: Chance of acceptance very low



every edge, chance of failure $\frac{\lambda^2}{(1+\lambda)^2}$

So for multiple edges (say $\propto$ ind. edges) $\left(1 - \frac{\lambda^2}{(1+\lambda)^2}\right)^{\infty}$ is upper bound on chance of success.



Solution PRAR

Recursive Acceptance Rejection

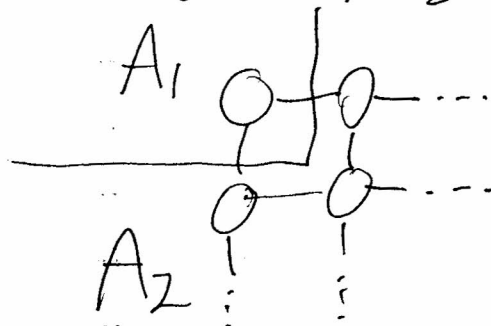
Repeat

Pick $v \in V$

Set $A_1 \leftarrow \{v\}$ w/ prob. $\lambda/(1+\lambda)$, otherwise $A_1 \leftarrow \{\}$

Draw A_2 from graph Π_λ on graph induced by $V - \{v\}$.

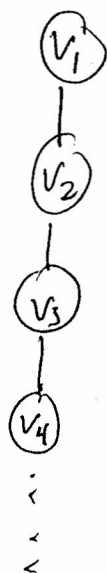
Until $A_1 \cup A_2$ is an independent set



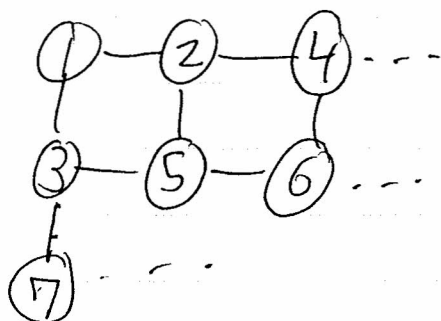
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Recursion Tree

Still Exponential



PRAR comes from simple idea that you don't need all recursion



Repeat

Pick any $v \in V$

~~Make~~ ^{Set} $A_i \leftarrow \{v\}$ w/ prob. $\frac{2}{1+\alpha}$, else set $A_i \leftarrow \{\}$.

If $A_i = \{v\}$

Let $N \leftarrow$ neighbors of v

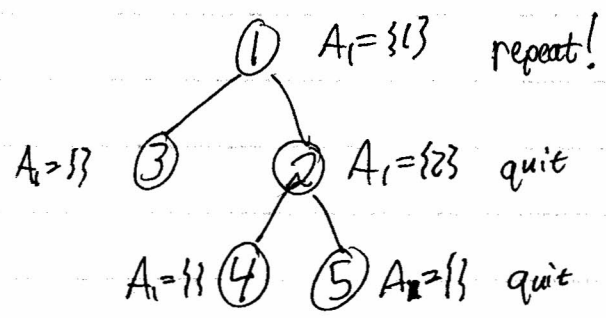
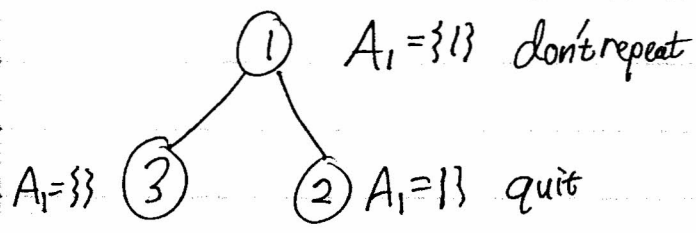
Recursion $\longrightarrow$ Find out whether nodes in N are part of ind. set.

Until $A_i = \{\}$ or no nodes in N part of ind. set.

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Typical Recursion Tree

① $A_1 = \{\}$ quit, don't repeat



Running time: Does tree grow or shrink on average?

T = mean time to resolve tree w/ one node

$$T \leq \underbrace{\frac{1}{1+\lambda}}_{\text{chance of } \{\}} \cdot (1) + \frac{\lambda}{1+\lambda} \left[\underbrace{T(\Delta-1)}_{\text{max degree of graph}} + \underbrace{T}_{\text{might reject and do original node over}} \right]$$

$$(1+\lambda)T \leq 1 + \lambda(\Delta-1)T$$

$$[1+\lambda - \lambda(\Delta-1)]T \leq 1$$

$$T \leq \frac{1}{1 - \lambda(\Delta-1)} \quad \text{for } 1 - \lambda(\Delta-1) \geq 0$$

$$\lambda \leq \frac{1}{\Delta-1}$$

[6]

PRAR not only method for generating independent sets

- Bounding chains [H '99]
- Randomness Recycler [Fill, H. 2001]

Which is better?

Bounding chains $\frac{2}{\Delta-2}$

RR $\frac{3/2}{\Delta-1}$

PRAR $\frac{1}{\Delta-1}$

Pessimist: Don't waste your time

Optimist: Opportunity for improvement

Conj.: Provably poly time for $\lambda < \frac{e}{\Delta-1}$