Fast approximation algorithms for partition functions of Gibbs distributions

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The Problem

Besag [1974] modeled soil plots as good (green) or bad (red)



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Basic inference



Finding $Z(\cdot)$

Need $Z(\cdot)$ to have posterior density

- Brute force computation long (curse of dimensionality)
- Brute force for Ising model takes 2^{# of vertices} time
- In denominator, so need relative (not additive) accuracy

Also need $Z(\cdot)$ for computing Bayes factors

Again, for Bayes factors relative accuracy essential

$Z(\beta)$: The Partition Function

Partition function of Gibbs distribution

Ingredients:

State space $\Omega \subseteq \mathbb{R}^n$ Nonnegative function H(x)Parameter β

Then the goal is to approximate:

$$Z(\beta) = \int_{x\in\Omega} \exp(-\beta H(x)) dx$$

Discrete version

Ingredients:

State space $\Omega \subseteq \{1, 2, ..., c\}^n$ Nonnegative function H(x)Parameter β

Then the goal is to approximate:

$$Z(\beta) = \sum_{x \in \Omega} \exp(-\beta H(x))$$

Terminology

Definition (Gibbs distribution)

X has a (discrete) *Gibbs distribution* π_{β} if for all $x \in \Omega$,

$$\mathbb{P}(X=x)=\frac{1}{Z(\beta)}\exp(-\beta H(x)).$$

Definition (Partition function)

The partition function of a Gibbs distribution is

$$Z(\beta) = \sum_{x \in \Omega} \exp(-\beta H(x)).$$

Today's result

An approximation algorithm where

- $ightharpoonup Given \epsilon > 0$
- Given the ability to sample from π'_{β} for $\beta' \in [0, \beta]$
- Outputs $\hat{Z}(\beta)$ so that

$$\mathbb{P}\left(rac{1}{1+\epsilon} \leq rac{\hat{Z}(eta)}{Z(eta)} \leq 1+\epsilon
ight) \geq 3/4$$

Let $q = \ln(Z(\beta))$. Method requires a number of samples

$$O(q \ln(M)(\ln(q) + \epsilon^{-2})), \quad M = \max_{x} H(x)$$

Is that good?

Why is problem hard?

Typically $Z(\beta)$ is exponential in *n*, the input size of problem

Many methods for lowering variance

- Multistage Sampling [Valleau Card 1972]
- Bridge Sampling [Meng Wong 1996]
- Nested Sampling [Skilling 2006]

The above are not approximation algorithms

No *guarantee* on quality of estimate obtained (But they could be faster in practice)

Approximation Algorithms

Let $q = \ln(Z(0)/Z(\beta))...$

Self-reducibility [Jerrum, Valiant, Vazirani 1986] $O[q^2 e^{-2}]$ time under best conditions

TPA [H., Schott 2010] $O[q^2 e^{-2}]$, simpler, constant ~ 20

SVV [Štefankovič, Vempala, Vigoda 2009] $O[q(\ln(q) + \ln(M))^5 e^{-2}]$ (constant ~ 10¹⁰)

Paired Product Approximation Algorithm [H. 2012] $O[q \ln(M)(e^{-2} + \ln(q))]$ (constant < 50)

The algorithm

New algorithm

Importance sampling

Multistage sampling

TPA

Paired Estimator

Importance Sampling

Usually Z(0) easy to find, so need to estimate

 $\frac{Z(\beta)}{Z(0)}$

For
$$X \sim \pi_0$$
, set $W = \exp(-\beta H(X))$
 $\mathbb{E}[W] = rac{Z(\beta)}{Z(0)}$

Relative variance (square of coefficient of variation)

$$\mathbb{V}_{\mathsf{rel}}(W) = rac{\mathbb{V}(W)}{\mathbb{E}[W]^2} = rac{Z(2eta)Z(0)}{Z(eta)} - 1$$

The relative variance picture





Only care about 0 to β , harder to sample at 2β !

Paired Estimator

Solution to 2β problem: use two estimators

$$X \sim \pi_0, \quad Y \sim \pi_\beta$$

$$W := \exp(-(\beta/2)H(X)), \quad V := \exp((\beta/2)H(Y))$$

W estimates
$$\frac{Z(\beta/2)}{Z(0)}$$
, *V* estimates $\frac{Z(\beta/2)}{Z(\beta)}$

The relative variance for paired estimator



$$\mathbb{V}_{\mathsf{rel}}(W) = \mathbb{V}_{\mathsf{rel}}(V) = \exp(2\delta)$$

But what to do about δ being too large?

Multistage sampling

Breaking the interval into pieces [Valleau and Card 1972]



$$\frac{Z(\beta)}{Z(0)} = \frac{Z(\beta_1)}{Z(0)} \frac{Z(\beta_2)}{Z(\beta_1)} \frac{Z(\beta)}{Z(\beta_2)}$$

Relative variance of product estimator

Relative variance of $\prod W_i$ (and $\prod V_i$)

$$-1 + \exp\left(\sum_i \delta_i\right)$$

New question:

How do we make $\sum_i \delta_i$ small without using too many intervals?

Making $\sum \delta_i$ small

Fact (Lemma 3.3 of H. 2012)

 $\sum \delta_i$ is small if the vertical drop in *z* over each interval is small.

Call such a set of intervals well balanced

- To get well balanced intervals use TPA [H. Schott 2010]
- Time to get/number of such intervals $O(q \ln(M) \ln(q))$
- Once have intervals, $O(q \ln(M)\epsilon^{-2})$ time to approximate $\prod W_i$ and $\prod V_i$

Well balanced schedules also good for Markov chains

 Necessary for fast parallel tempering (see [Woodard, Schmidler, H. 2009])

The Algorithm

- Obtain an initial estimate of q with TPA
- Obtain a well balanced set of intervals with TPA
- **③** Use the well balanced schedule to get \bar{W} , \bar{V}
- Final estimate \bar{W}/\bar{V}

Conclusions

Most important thing

Previous practical methods quadratic in q (roughly the dimension), new method nearly linear

Uses some old ideas

Importance Sampling, Multistage Sampling, Product Estimator

Add some new ideas

Paired Product Estimator, Well-balanced schedule through TPA

The resulting approximation algorithm

Simple to implement, guaranteed quality of estimate

References

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